# CSCE 658: Randomized Algorithms

Lecture 21

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# Relevant Supplementary Material

 Chapter 3-4 of "The Algorithmic Foundations of Differential Privacy", by Cynthia Dwork and Aaron Roth (https://www.cis.upenn.edu/~aaroth/Papers/privacybook.pd f)

#### Last Time: Differential Privacy

• [DMNS06] Given  $\varepsilon > 0$  and  $\delta \in (0,1)$ , a randomized algorithm  $A: U^* \to Y$  is  $(\varepsilon, \delta)$ -differentially private if, for every neighboring pair D and D' of datasets, and for all  $E \subseteq Y$ ,  $\Pr[A(D) \in E] \le e^{\varepsilon} \cdot \Pr[A(D') \in E] + \delta$ 



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• Implication: Deterministic algorithms cannot be differentially private unless they are a constant function

# Last Time: Local Differential Privacy (LDP)

• [KLNRS08] Given  $\varepsilon > 0$  and  $\delta \in (0,1)$ , a randomized algorithm  $A: U^* \to Y$  is  $(\varepsilon, \delta)$ -differentially private if, for every pairs of users' possible data x and x' and for all  $E \subseteq Y$ ,  $\Pr[A(x) \in E] \leq e^{\varepsilon} \cdot \Pr[A(x') \in E] + \delta$ 

- Algorithm takes a single user's data
- Compared to previous definition of DP, where algorithm takes all users' data

# Last Time: Randomized Response, Revisisted

- How many people in this class have a pet?
- Generate a random integer:
  - If it is even, answer truthfully
  - Otherwise, proceed below
- Generate another random integer:
  - If it is even, answer YES
  - Otherwise if it is odd, answer NO



# Last Time: Randomized Response, Revisisted

- Answer is correct in expectation, but what is its variance?
- Each answer is incorrect with probability  $\frac{1}{4}$
- The variance is O(n)

- By Chebyshev, we have additive error  $O(\sqrt{n})$  with probability 0.99
- By anti-concentration, error is  $\Omega(\sqrt{n})$



#### Last Time: Laplace Mechanism

- Goal: Algorithm computes f(x) and releases f(x) + Z, where  $Z \sim$  $Lap\left(\frac{\sigma_f}{\varepsilon}\right)$
- Laplacian distribution: Probability density function for Lap(b) is  $p(x) = \frac{1}{2b} \exp\left(-\frac{|x|}{b}\right) = \frac{1}{2b} e^{\left(-\frac{|x|}{b}\right)}$



#### Laplace Mechanism

 Theorem: Laplace mechanism is *ɛ*-differentially private (pure DP)

- What is the algorithm for the private counting problem by the Laplace mechanism?
- What is the error for the private counting problem by the Laplace mechanism?

• How do we answer non-numeric queries, e.g., selection?

• Example: What is the most common eye color in the room?

 Example: Suppose a study is conducted that finds the current location of individuals, in the two-dimensional plane

• Who is the closest individual to a query location?

• What if we want to output the "best" answer, but noise can significantly destroy the answer?

- Example: Suppose we have a large number of apples, and A, B, C each bid \$1.00 and D bids \$4.01. What is the optimal price?
- At \$4.01 the revenue, the revenue is \$4.01, at \$4.00 and at \$1.00 the revenue is \$4.00, but at \$3.02 the revenue is zero!

• Choose a score function  $S: (X^n, Y) \rightarrow \mathbb{R}$  and global sensitivity  $\sigma$ 

• Sample  $y \in Y$  with probability proportional to  $\exp\left(\frac{\varepsilon}{2\sigma}S(x,y)\right)$ 

• How do we answer non-numeric queries, e.g., selection?

• Example: What is the most common eye color in the room?

 Example: Suppose a study is conducted that finds the current location of individuals, in the two-dimensional plane

• Who is the closest individual to a query location?

- Theorem: Exponential mechanism is *e*-differentially private (pure DP)
- Suppose the "answer" for dataset *D* is *z* and the "answer" for dataset *D* is *z*'

• Let y be the output of the exponential mechanism for Dand let y' be the output of the exponential mechanism for D'

• Want to show that for any fixed x,  $\frac{\Pr[y=x]}{\Pr[y'=x]} \le e^{\varepsilon}$ 

$$\frac{\Pr[y=x]}{\Pr[y'=x]} \le \exp\left(\frac{\varepsilon S(D,x)}{2\sigma_f}\right) / \sum_{x} \exp\left(\frac{\varepsilon S(D,x)}{2\sigma_f}\right)$$
$$\cdot \sum_{x} \exp\left(\frac{\varepsilon S(D',x)}{2\sigma_f}\right) / \exp\left(\frac{\varepsilon S(D',x)}{2\sigma_f}\right)$$

• Want to show that for any fixed x,  $\frac{\Pr[y=x]}{\Pr[y'=x]} \leq e^{\varepsilon}$ 

$$\frac{\Pr[y=x]}{\Pr[y'=x]} \le \exp\left(\frac{\varepsilon(S(D,x) - S(D',x))}{2\sigma_f}\right)$$
$$\cdot \sum_{x} \exp\left(\frac{\varepsilon S(D,x)}{2\sigma_f}\right) / \sum_{x} \exp\left(\frac{\varepsilon S(D',x)}{2\sigma_f}\right)$$

• Want to show that for any fixed x,  $\frac{\Pr[y=x]}{\Pr[y'=x]} \le e^{\varepsilon}$ 

 $\frac{\Pr[y = x]}{\Pr[y' = x]} \le \exp\left(\frac{\varepsilon}{2}\right)$  $\cdot \exp\left(\frac{\varepsilon}{2}\right) \sum_{x} \exp\left(\frac{\varepsilon S(D, x)}{2\sigma_f}\right) / \sum_{x} \exp\left(\frac{\varepsilon S(D, x)}{2\sigma_f}\right)$  $= \exp(\varepsilon)$ 

- Theorem: Exponential mechanism is *e*-differentially private (pure DP)
- Note we can still apply exponential mechanism when Y is the set of the real numbers

• How does it compare to the Laplace mechanism?

- Consider a query with sensitivity  $\sigma_f$
- Suppose the "answer" for dataset D is z and the "answer" for dataset D is z'

• Let 
$$y = z + \operatorname{Lap}\left(\frac{\sigma_f}{\varepsilon}\right)$$
  
•  $p(x) = \frac{\varepsilon}{2\sigma_f} \exp\left(-\frac{\varepsilon|x|}{\sigma_f}\right) = \frac{\varepsilon}{2\sigma_f} e^{\left(-\frac{\varepsilon|x|}{\sigma_f}\right)}$  for a Laplace distribution with scale parameter  $\frac{\sigma_f}{\varepsilon}$ 

• Laplace mechanism: Output y = z + x with probability proportional to  $\frac{\varepsilon}{2\sigma_f} e^{\left(-\frac{\varepsilon|x|}{\sigma_f}\right)}$ 

- Consider a query with sensitivity  $\sigma_f$
- Suppose the "answer" for dataset *D* is *z* and the "answer" for dataset *D* is *z*'
- Choose score function -2|y-x|
- Exponential mechanism: Output y = z + x with probability proportional to  $e^{\left(-\frac{\varepsilon|x|}{\sigma_f}\right)}$

- Laplace mechanism: Output y = z + x with probability proportional to  $\frac{\varepsilon}{2\sigma_f} e^{\left(-\frac{\varepsilon|x|}{\sigma_f}\right)}$
- Exponential mechanism: Output y = z + x with probability proportional to  $e^{\left(-\frac{\varepsilon|x|}{\sigma_f}\right)}$
- Recovers the Laplace mechanism!

## Exponential Mechanism Drawbacks

• Sampling process may be inefficient

• Error can be large