CSCE 658: Randomized Algorithms

Lecture 3

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Class Logistic Updates

• Course materials: <u>https://samsonzhou.github.io/CSCE658-</u> <u>S24/csce658-s24</u>

Class Logistic Updates

- Problem Set 1 posted, due next Thursday, February 1, 2024, 5 pm CT
- PS1 groups sent via e-mail, please confirm receipt by replyall

- Submit PS1 via e-mail as a PDF, typeset in LaTeX
- LaTeX template for PS1 available on class webpage, for your convenience

Trivia Question #1 (Birthday Paradox)

- Suppose we have a fair *n*-sided die. How many times should we roll the die before the probability we see a repeated outcome among the rolls is at least $\frac{1}{2}$? Example: 1, 5, 2, 4, 5
- O(1)
- $\Theta(\log n)$
- $\Theta(\sqrt{n})$
- **Θ**(*n*)

Trivia Question #2 (Limits)

• Let
$$c > 0$$
 be a constant. What is $\lim_{n \to \infty} \left(1 - \frac{c}{n}\right)^n$?

• 0
•
$$\frac{1}{c}$$

• $\frac{1}{2c}$
• $\frac{1}{e^{c}}$
• 1

• Suppose we have a room with 367 people. What is the probability that two people share the same birthday?

• Suppose we have a room with 367 people. What is the probability that two people share the same birthday?

• Suppose we have a room with 23 people. What is the probability that two people share the same birthday?

$$\left(1-\frac{0}{n}\right)$$

$$\left(1-\frac{0}{n}\right)\left(1-\frac{1}{n}\right)$$

$$\left(1-\frac{0}{n}\right)\left(1-\frac{1}{n}\right)\left(1-\frac{2}{n}\right)$$

$$\left(1-\frac{0}{n}\right)\left(1-\frac{1}{n}\right)\dots\left(1-\frac{k-1}{n}\right)$$

$$\left(1-\frac{0}{n}\right)\left(1-\frac{1}{n}\right)\dots\left(1-\frac{k-1}{n}\right) < \frac{1}{2}$$

$$\left(1-\frac{0}{n}\right)\left(1-\frac{1}{n}\right)\dots\left(1-\frac{k-1}{n}\right) < \frac{1}{2} \qquad \text{for} \quad k = O(\sqrt{n})$$

- Suppose we have a fair *n*-sided die. "On average", how many times should we roll the die before we see a repeated outcome among the rolls?
- $O(\sqrt{n})$
- But is it $\Theta(\sqrt{n})$?

- Let S_i be the event that the *i*-th roll is a repeated outcome, conditioned on the previous rolls not being a repeated outcome
- $\Pr[S_i] = \frac{i-1}{n}$
- $\Pr[S_1 \cup \cdots \cup S_k] \le ???$

- Let S_i be the event that the *i*-th roll is a repeated outcome, conditioned on the previous rolls not being a repeated outcome
- $\Pr[S_i] = \frac{i-1}{n}$
- $\Pr[S_1 \cup \dots \cup S_k] \le \frac{0}{n} + \dots + \frac{k-1}{n} \le \frac{k^2}{n}$

 Suppose we have a fair *n*-sided die that we roll *k* = 1, 2, 3, 4,... times. What is the probability we DO NOT see a repeated outcome among the rolls?

- Let S_i be the event that the *i*-th roll is a repeated outcome, conditioned on the previous rolls not being a repeated outcome
- $\Pr[S_i] = \frac{i-1}{n}$

•
$$\Pr[S_1 \cup \dots \cup S_k] \le \frac{0}{n} + \dots + \frac{k-1}{n} \le \frac{k^2}{n}$$

Union Bound

- Suppose we have a fair *n*-sided die. "On average", how many times should we roll the die before we see a repeated outcome among the rolls?
- $\Theta(\sqrt{n})$

- We are trying to learn a new language on an app, which claims to have a database of *1 million words*
- Each time we ask the app, it gives us a random word in the database
- We want to verify the claim





- We could use the app until we see 1 million unique words, but that would take at least 1 million checks
- Instead, we use the app for 1000 times and count the number of pairwise duplicates
- If there are many duplicates, the database is probably not very large





- We use the app for *k* times and count the number of pairwise duplicates
- If we see the same word on the 3-rd time, the 100-th time, and the 205-th time, there are 3 pairwise duplicates: (3, 100), (3, 205), (100, 205)





Expected Value

• The expected value of a random variable X over a sample space Ω is:

$$\mathbf{E}[X] = \sum_{x \in \Omega} \Pr[X = x] \cdot x$$

• The "average value of the random variable"

Expected Value

- Suppose we roll a 6-sided die
- Let *X* be the outcome of the roll
- What is **E**[X]?

$$\mathbf{E}[X+Y] = \sum_{x \in \Omega_X} \sum_{y \in \Omega_Y} \Pr[X = x, Y = y] \cdot (x+y)$$

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=
$$\sum_{x \in \Omega_X} \sum_{y \in \Omega_Y} \Pr[X = x, Y = y] \cdot x + \sum_{x \in \Omega_X} \sum_{y \in \Omega_Y} \Pr[X = x, Y = y] \cdot y$$

$$\mathbf{E}[X+Y] = \sum_{x \in \Omega_{X}} \sum_{y \in \Omega_{Y}} \Pr[X = x, Y = y] \cdot (x+y)$$

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=
$$\sum_{x \in \Omega_{X}} x \sum_{y \in \Omega_{Y}} \Pr[X = x, Y = y] + \sum_{y \in \Omega_{Y}} y \sum_{x \in \Omega_{X}} \Pr[X = x, Y = y]$$

$$\mathbf{E}[X+Y] = \sum_{x \in \Omega_X} \sum_{y \in \Omega_Y} \Pr[X = x, Y = y] \cdot (x+y)$$

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=
$$\sum_{x \in \Omega_X} x \sum_{y \in \Omega_Y} \Pr[X = x, Y = y] + \sum_{y \in \Omega_Y} y \sum_{x \in \Omega_X} \Pr[X = x, Y = y]$$

=
$$\sum_{x \in \Omega_X} x \cdot \Pr[X = x] + \sum_{y \in \Omega_Y} y \cdot \Pr[Y = y] = \mathbf{E}[X] + \mathbf{E}[Y]$$

$$\left(1-\frac{0}{n}\right)\left(1-\frac{1}{n}\right)\dots\left(1-\frac{k-1}{n}\right)$$

• Suppose we have a fair *n*-sided die that we roll k = 1, 2, 3, 4,... times. What is the expected number of pairwise collisions among the rolls?

• Let X_i be the number of pairwise collisions on the *i*-th roll

• We have
$$E[X_i] = \frac{i-1}{n}$$

• Let X be the number of pairwise collisions after k rolls

• What is **E**[X]?

• Let X be the number of pairwise collisions after k rolls

 $E[X] = E[X_1 + \dots + X_k]$ $= E[X_1] + \dots + E[X_k]$ $= \frac{0}{n} + \dots + \frac{k-1}{n}$ $= \frac{k(k-1)}{2n}$

• $\operatorname{E}[X] = \frac{k(k-1)}{2n}$

•
$$\frac{(k-1)^2}{2n} \le \operatorname{E}[X] \le \frac{k^2}{2n}$$

• $k = 2\sqrt{n} + 1$ implies $E[X] \ge 1$

•
$$k = \frac{\sqrt{n}}{2}$$
 implies $\mathbb{E}[X] \le \frac{1}{4}$

- We use the app for k = 1000 times and count the number of pairwise duplicates
- If the database contains 1 million words, the expected number of pairwise duplicates is $E[X] = \frac{k(k-1)}{2n} < 0.5$





- If the database contains 1 million words, the expected number of pairwise duplicates is $E[X] = \frac{k(k-1)}{2n} < 0.5$
- ...We see 20 duplicates
- We think the claim is incorrect, but how can we be sure?





Concentration Inequalities

 Concentration inequalities bound the probability that a random variable is "far away" from its expectation

 Often used in understanding the performance of statistical tests, the behavior of data sampled from various distributions, and for our purposes, the guarantees of randomized algorithms

Markov's Inequality

• Let $X \ge 0$ be a non-negative random variable. Then for any t > 0:

$$\Pr[X \ge t \cdot E[X]] \le \frac{1}{t}$$

Proof of Markov's Inequality

• Let $X \ge 0$ be a non-negative random variable. Then for any t > 0:

 $E[X] = \sum_{x \in \Omega} \Pr[X = x] \cdot x$ $= \sum_{x \ge t \cdot E[X]} \Pr[X = x] \cdot x + \sum_{x < t \cdot E[X]} \Pr[X = x] \cdot x$ $\ge \sum_{x \ge t \cdot E[X]} \Pr[X = x] \cdot x$ $\ge t \cdot E[X] \sum_{x \ge t \cdot E[X]} \Pr[X = x]$ $= t \cdot E[X] \cdot \Pr[X \ge t \cdot E[X]]$

$$\left(1-\frac{0}{n}\right)\left(1-\frac{1}{n}\right)\dots\left(1-\frac{k-1}{n}\right)$$

• Suppose we have a fair *n*-sided die that we roll k = 1, 2, 3, 4,... times. What is the expected number of pairwise collisions among the rolls?

• Let X_i be the number of pairwise collisions on the *i*-th roll

• We have
$$E[X_i] = \frac{i-1}{n}$$

• $\operatorname{E}[X] = \frac{k(k-1)}{2n}$

•
$$\frac{(k-1)^2}{2n} \le \operatorname{E}[X] \le \frac{k^2}{2n}$$

• $k = 2\sqrt{n} + 1$ implies $E[X] \ge 1$

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$$k = \frac{\sqrt{n}}{2}$$
 implies $\mathbb{E}[X] \le \frac{1}{4}$

•
$$\mathbf{E}[X] = \frac{k(k-1)}{2n}$$

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$$\frac{(k-1)^2}{2n} \le \operatorname{E}[X] \le \frac{k^2}{2n}$$

• $k = 2\sqrt{n} + 1$ implies $E[X] \ge 1$

•
$$k = \frac{\sqrt{n}}{2}$$
 implies $E[X] \le \frac{1}{4}$, and by Markov's inequality, $\Pr[X \ge 1] \le \frac{1}{4}$

- If the database contains 1 million words, the expected number of pairwise duplicates is $E[X] = \frac{k(k-1)}{2n} < 0.5$
- ...We see 20 duplicates
- We think the claim is incorrect, but how can we be sure?





- If the database contains 1 million words, the expected number of pairwise duplicates is $E[X] = \frac{k(k-1)}{2n} < 0.5$
- ...We see 20 duplicates
- $\Pr[X \ge 20] \le \frac{1}{40}$



