CSCE 689: Special Topics in Modern Algorithms for Data Science

Lecture 17

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Previously: Variance

• The variance of a random variable X over Ω is:

$$Var[X] = E[X^2] - (E[X])^2$$

• Linearity of variance for *independent* random variables: Var[X + Y] = Var[X] + Var[Y]

Previously: Chebyshev's Inequality

• Let X be a random variable with expected value $\mu \coloneqq E[X]$ and variance $\sigma^2 \coloneqq Var[X]$

•
$$\Pr[|X - E[X]| \ge t] \le \frac{\operatorname{Var}[X]}{t^2}$$
 becomes $\Pr[|X - E[X]| \ge t] \le \frac{\sigma^2}{t^2}$

$$\Pr[|X - \mu| \ge k\sigma] \le \frac{1}{k^2}$$

• "Bounding the deviation of a random variable in terms of its variance"

Last Time: Distinct Elements (F_0 Estimation)

- Given a set S of m elements from [n], let f_i be the frequency of element i. (How often it appears)
- Let F_0 be the frequency moment of the vector:

$$F_0 = |\{i : f_i \neq 0\}|$$

• Goal: Given a set S of m elements from [n] and an accuracy parameter ε , output a $(1 + \varepsilon)$ -approximation to F_0













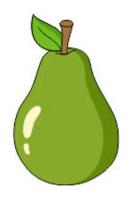












• Intuition: How is this done in practice?

• Let S be a set of N numbers

• Suppose we form set S' by sampling each item of S with probability $\frac{1}{2}$

• How many numbers are in S'?

• Let S be a set of N numbers

• Suppose we form set S' by sampling each item of S with probability $\frac{1}{2}$

• Can we use S' to get a good estimate of N?

• Let S be a set of N numbers, suppose we form set S' by sampling each item of S with probability $\frac{1}{2}$

• We have $E[|S'|] = \frac{N}{2}$ and $Var[|S'|] \le \frac{N}{2}$

• Claim: We have $Var[|S'|] \le \frac{N}{2}$

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- $Var[X_i] = E[X_i^2] E[X_i]^2$

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- $Var[|S'|] = \frac{N}{4}$

• What can we say about $\Pr\left[\left||S'| - \frac{N}{2}\right| \ge t\right]$?

• By Chebyshev's inequality, we have $\Pr\left[\left||S'| - \frac{N}{2}\right| \ge 100\sqrt{N}\right] \le \frac{1}{10}$

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$$\frac{N}{2} - 100\sqrt{N} \le |S'| \le \frac{N}{2} + 100\sqrt{N}$$

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• If $200\sqrt{N} \le \frac{N}{100}$, then $N - 200\sqrt{N} \le 2|S'| \le N + 200\sqrt{N}$ implies

$$0.99N \le 2|S'| \le 1.01N$$

Very good approximation to N

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- Sample each item of the *universe* with probability $\frac{1}{2}$, acquire new universe U'
- Let S' be the items in the data stream that are in U'
- Output 2|*S'*|

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What's the problem with this approach?

• Let S be a set of N numbers

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• Suppose we form set S' by sampling each item of S with probability p

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• We have $E[|S'|] = \frac{N}{2}$ and $Var[|S'|] \le \frac{N}{2}$

• Let S be a set of N numbers, suppose we form set S' by sampling each item of S with probability p

• We have E[|S'|] = pN and $Var[|S'|] \le pN$

• (S' is formed by sampling each item of S with probability $\frac{1}{2}$) With probability at least $\frac{9}{10}$,

$$\frac{N}{2} - 100\sqrt{N} \le |S'| \le \frac{N}{2} + 100\sqrt{N}$$

• Thus with probability at least $\frac{9}{10}$,

$$N - 200\sqrt{N} \le 2|S'| \le N + 200\sqrt{N}$$

• (S' is formed by sampling each item of S with probability p) With probability at least $\frac{9}{10}$,

$$pN - 100\sqrt{pN} \le |S'| \le pN + 100\sqrt{pN}$$

• Thus with probability at least $\frac{9}{10}$,

$$N - \frac{100}{\sqrt{p}} \sqrt{N} \le \frac{1}{p} |S'| \le N + \frac{100}{\sqrt{p}} \sqrt{N}$$

• (S' is formed by sampling each item of S with probability p) With probability at least $\frac{9}{10}$,

$$N - \frac{100}{\sqrt{p}} \sqrt{N} \le \frac{1}{p} |S'| \le N + \frac{100}{\sqrt{p}} \sqrt{N}$$

• If
$$\frac{100}{\sqrt{p}}\sqrt{N} \le \varepsilon N$$
, then $N - \frac{100}{\sqrt{p}}\sqrt{N} \le \frac{1}{p}|S'| \le N + \frac{100}{\sqrt{p}}\sqrt{N}$ implies

$$(1-\varepsilon)N \le \frac{1}{p}|S'| \le (1+\varepsilon)N$$

• In other words, with probability at least $\frac{9}{10}$, we have that $\frac{1}{p}|S'|$ is a $(1+\varepsilon)$ -approximation of N

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- Recall, we required $\frac{100}{\sqrt{p}}\sqrt{N} \le \varepsilon N$, so $p \ge \frac{1000}{\varepsilon^2 N}$

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- What is the problem here?

Must know N to set p, but the goal is to find N!

• Observation: We do not need $p=\frac{1000}{\varepsilon^2 N}$, it is also fine to have $p=\frac{2000}{\varepsilon^2 N}$

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• Not too many samples, i.e., S' is small, but enough to find a good approximation to N

- For $p = \Theta\left(\frac{1}{\varepsilon^2 N}\right)$:
 - p is large enough to find a good approximation to N
 - We have $E[|S'|] = pN = \Theta\left(\frac{1}{\varepsilon^2}\right)$

• We want p such that $\mathrm{E}[|S'|] = pN = \Theta\left(\frac{1}{\varepsilon^2}\right)$

• Intuition: Try $p=1,\frac{1}{2},\frac{1}{4},\frac{1}{8},\frac{1}{16},\ldots$, and see which one has

$$\frac{1000}{\varepsilon^2} \le |S'| \le \frac{2000}{\varepsilon^2}$$

• With high probability, one of these guesses will have $\frac{1000}{\varepsilon^2} \le |S'| \le \frac{2000}{\varepsilon^2}$

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However, the wrong guesses will have too many samples

• Fix: Dynamically changing guess for p and subsampling

• Algorithm: Set $U_0 = [n]$ and for each i, sample each element of U_{i-1} into U_i with probability $\frac{1}{2}$

• Start index i=0 and track the number $|S\cap U_i|$ of elements of S in U_i

• If $|S \cap U_i| > \frac{2000}{\varepsilon^2} \log n$, then increment i = i + 1

• At the end of the stream, output $2^i \cdot |S \cap U_i|$

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• If
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, then increment $i = i + 1 \left(\frac{1}{p}\right)$

• At the end of the stream, output $2^i \cdot |S \cap U_i|$

• Recall that $\frac{1}{p}|S'|$ is a $(1 + \varepsilon)$ -approximation of N

• $2^i \cdot |S \cap U_i|$ is a $(1 + \varepsilon)$ -approximation of N

• At the end of the stream, output $2^i \cdot |S \cap U_i|$

• Summary: Algorithm stores at most $\frac{2000}{\varepsilon^2} \log n$ elements from the stream, uses $\Theta\left(\frac{1}{\varepsilon^2} \log n\right)$ words of space