CSCE 689: Special Topics in Modern Algorithms for Data Science

Lecture 2

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Last Time: Class Logistics

- Course materials: https://samsonzhou.github.io/csce689-2023
- LaTeX summary of lectures 20%
- Midterm presentation 35%
- Final project 45%

Last Time: Probability Basics

• Conditional distribution: $\Pr[X = x | Y = y]$ is the probability that X achieves the value x when Y achieves the value y

$$\Pr[X = x | Y = y] = \frac{\Pr[X = x, Y = y]}{\Pr[Y = y]}$$

- Implies Bayes' theorem
- Random variables X and Y are independent if $\Pr[X = x] = \Pr[X = x | Y = y]$ for all possible outcomes $x \in \Omega_X$, $y \in \Omega_Y$

Warm-Up Question

• Suppose S_1 is a "bad" event that occurs with probability $\frac{0}{n}$ • Suppose S_2 is a "bad" event that occurs with probability $\frac{1}{n}$ • Suppose S_3 is a "bad" event that occurs with probability $\frac{2}{n}$

• What is the probability that none of the bad events occurs?

Warm-Up Question

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• What is *a lower bound* on the probability that none of the bad events occur?

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• What is *a lower bound* on the probability that none of the bad events occur? $1 - \frac{3}{n}$

Last Time: Union Bound (Boole's Inequality)

• Let $S_1, ..., S_k$ be a set of events that occur with probability $p_1, ..., p_k$

• The probability that at least one of the events $S_1, ..., S_k$ occurs is at most $p_1 + \cdots + p_k$

• Implication: the probability that NONE of the events $S_1, ..., S_k$ occur is at least $1 - (p_1 + \cdots + p_k)$

Last Time: Union Bound

• $\Pr[A \cup B] = \Pr[A] + \Pr[B] - \Pr[A \cap B]$



• Proof by induction

Today

- Hashing
- Abstraction: balls-in-bins
- Birthday paradox

Trivia Question #1 (Birthday Paradox)

- Suppose we have a fair *n*-sided die. "On average", how many times should we roll the die before we see a repeated outcome among the rolls? Example: 1, 5, 2, 4, 5
- $\Theta(1)$
- $\Theta(\log n)$
- $\Theta(\sqrt{n})$
- **Θ**(*n*)

Trivia Question #2 (Limits)

• Let
$$c > 0$$
 be a constant. What is $\lim_{n \to \infty} \left(1 - \frac{c}{n}\right)^n$?

• 0
•
$$\frac{1}{c}$$

• $\frac{1}{2c}$
• $\frac{1}{e^{c}}$
• 1













• Suppose we have a number of files, how do we consistently store them in memory?



• Goal: Fast query time





- We have a set of *m* items from some large universe that we want to store into a database (images, text documents, IP addresses) with *n* locations
- Goal: query(x) to check if the database contains x in O(1) time
- Hash function $h: U \rightarrow [n]$ maps items from the universe to a location in the database

Collisions

- Hash function $h: U \rightarrow [n]$ maps items from the universe to a location in the database
- For $|U| \gg n$, many items map to the same location
- Collision: when multiple items should be stored in the same location



Dealing with Collisions

- Many ways of dealing with collisions
 - Store multiple items in the same location as a linked list
 - Bump item to the next available spot
 - Bump item to the next available spot using another hash function
 - Power-of-two-choices

Dealing with Collisions

• Suppose we store multiple items in the same location as a linked list



• If the maximum number of collisions in a location is *c*, then could traverse a linked list of size *c* for a query

• Query runtime: O(c)

Dealing with Collisions

• Goal: minimize *c*, the maximum number of collisions in a location

• In the worst case, all items could hash to the same location, c = m

• Assume the hash function *h* is chosen "randomly"

Random Hash Function

- Let $h: U \to [n]$ be a random hash function, so that for each $x \in U$, we have that $\Pr[h(x) = i] = \frac{1}{n'}$, for all $i \in [n]$
- Assume independence, i.e., h(x) and h(y) are independent for any $x, y \in U$
- Suppose we insert *m* elements into a hash table with *n* locations using a random hash function. How do we analyze the number of pairwise collisions?

• Suppose we have a room with 367 people. What is the probability that two people share the same birthday?

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• Suppose we have a room with 23 people. What is the probability that two people share the same birthday?

$$\left(1-\frac{0}{n}\right)$$

$$\left(1-\frac{0}{n}\right)\left(1-\frac{1}{n}\right)$$

$$\left(1-\frac{0}{n}\right)\left(1-\frac{1}{n}\right)\left(1-\frac{2}{n}\right)$$

$$\left(1-\frac{0}{n}\right)\left(1-\frac{1}{n}\right)\dots\left(1-\frac{k-1}{n}\right)$$

$$\left(1-\frac{0}{n}\right)\left(1-\frac{1}{n}\right)\dots\left(1-\frac{k-1}{n}\right) < \frac{1}{2} \qquad \text{for} \quad k = O(\sqrt{n})$$

- Suppose we have a fair *n*-sided die. "On average", how many times should we roll the die before we see a repeated outcome among the rolls?
- $O(\sqrt{n})$
- But is it $\Theta(\sqrt{n})$?

- Let S_i be the event that the *i*-th roll is a repeated outcome, conditioned on the previous rolls not being a repeated outcome
- $\Pr[S_i] = \frac{i-1}{n}$
- $\Pr[S_1 \cup \cdots \cup S_k] \le ???$

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- $\Pr[S_i] = \frac{i-1}{n}$
- $\Pr[S_1 \cup \dots \cup S_k] \le \frac{0}{n} + \dots + \frac{k-1}{n} \le \frac{k^2}{n}$

 Suppose we have a fair *n*-sided die that we roll *k* = 1, 2, 3, 4,... times. What is the probability we DO NOT see a repeated outcome among the rolls?

- Let S_i be the event that the *i*-th roll is a repeated outcome, conditioned on the previous rolls not being a repeated outcome
- $\Pr[S_i] = \frac{i-1}{n}$

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$$\Pr[S_1 \cup \dots \cup S_k] \le \frac{0}{n} + \dots + \frac{k-1}{n} \le \frac{k^2}{n}$$

Union Bound

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- $\Theta(\sqrt{n})$

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