

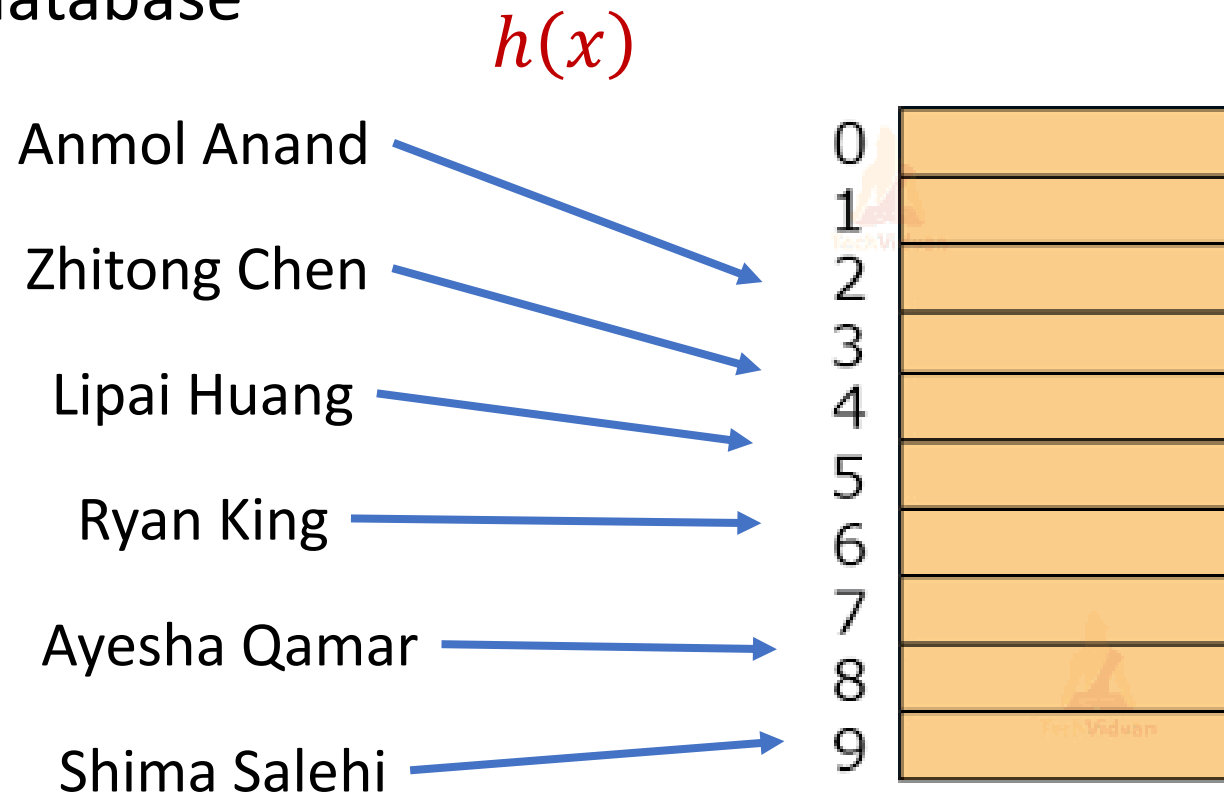
CSCSE 689: Special Topics in Modern Algorithms for Data Science

Lecture 3

Samson Zhou

Last Time: Hashing

- Hashing is a method to quickly map items from a universe to a location in a database



Last Time: Birthday Paradox

- Suppose we have a fair n -sided die. “On average”, how many times should we roll the die before we see a repeated outcome among the rolls? Example: 1, 5, 2, 4, 5
- $\Theta(1)$
- $\Theta(\log n)$
- $\Theta(\sqrt{n})$
- $\Theta(n)$

Last Time: Birthday Paradox

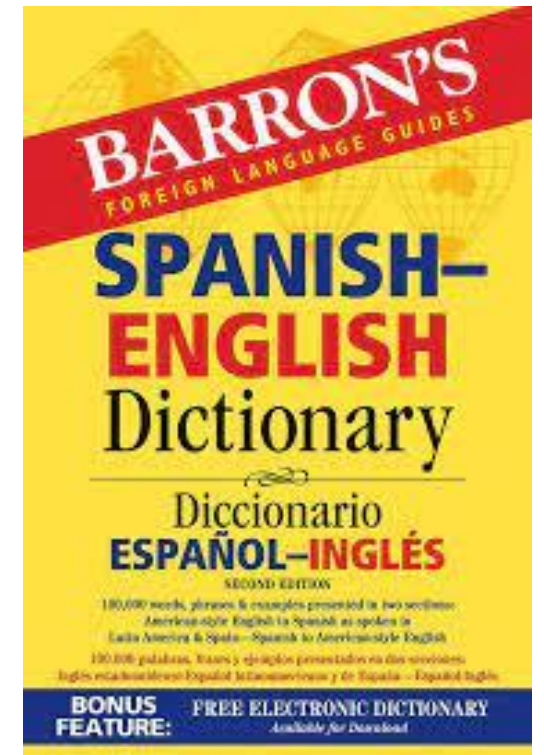
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Future

- **Next Monday:** Sign up for LaTeX scribe note slots
- **Today:** Meet your classmates (1)
- **Next Monday:** Meet your classmates (2), receive and consider list of potential projects/groups
- **Next Wednesday:** Discuss potential project groups
- **Next Friday:** Email me the members/group name

Case Study

- We could use the app until we see 1 million unique words, but that would take at least *1 million checks*
- Instead, we use the app for *1000 times* and count the number of pairwise duplicates
- If there are many duplicates, the database is probably not very large



Expected Value

- The expected value of a random variable X over Ω is:

$$E[X] = \sum_{x \in \Omega} \Pr[X = x] \cdot x$$

- The “average value of the random variable”
- Linearity of expectation: $E[X + Y] = E[X] + E[Y]$

Expected Value

- Suppose we roll a 6-sided die
- Let X be the outcome of the roll
- What is $E[X]$?

Linearity of Expectation

- Linearity of expectation: $E[X + Y] = E[X] + E[Y]$

$$E[X + Y] = \sum_{x \in \Omega_X} \sum_{y \in \Omega_Y} \Pr[X = x, Y = y] \cdot (x + y)$$

Linearity of Expectation

- Linearity of expectation: $E[X + Y] = E[X] + E[Y]$

$$\begin{aligned} E[X + Y] &= \sum_{x \in \Omega_X} \sum_{y \in \Omega_Y} \Pr[X = x, Y = y] \cdot (x + y) \\ &= \sum_{x \in \Omega_X} \sum_{y \in \Omega_Y} \Pr[X = x, Y = y] \cdot x + \sum_{x \in \Omega_X} \sum_{y \in \Omega_Y} \Pr[X = x, Y = y] \cdot y \end{aligned}$$

Linearity of Expectation

- Linearity of expectation: $E[X + Y] = E[X] + E[Y]$

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Linearity of Expectation

- Linearity of expectation: $E[X + Y] = E[X] + E[Y]$

$$\begin{aligned} E[X + Y] &= \sum_{x \in \Omega_X} \sum_{y \in \Omega_Y} \Pr[X = x, Y = y] \cdot (x + y) \\ &= \sum_{x \in \Omega_X} \sum_{y \in \Omega_Y} \Pr[X = x, Y = y] \cdot x + \sum_{x \in \Omega_X} \sum_{y \in \Omega_Y} \Pr[X = x, Y = y] \cdot y \\ &= \sum_{x \in \Omega_X} x \sum_{y \in \Omega_Y} \Pr[X = x, Y = y] + \sum_{y \in \Omega_Y} y \sum_{x \in \Omega_X} \Pr[X = x, Y = y] \\ &= \sum_{x \in \Omega_X} x \cdot \Pr[X = x] + \sum_{y \in \Omega_Y} y \cdot \Pr[Y = y] = E[X] + E[Y] \end{aligned}$$

Birthday Paradox

- Suppose we have a fair n -sided die that we roll $k = 1, 2, 3, 4, \dots$ times. What is the probability we **DO NOT** see a repeated outcome among the rolls?

$$\left(1 - \frac{0}{n}\right) \left(1 - \frac{1}{n}\right) \dots \left(1 - \frac{k-1}{n}\right)$$

Birthday Paradox, Revisited

- Suppose we have a fair n -sided die that we roll $k = 1, 2, 3, 4, \dots$ times. What is the expected number of pairwise collisions among the rolls?
- Let X_i be the number of pairwise collisions on the i -th roll
- We have $E[X_i] = \frac{i-1}{n}$

Birthday Paradox, Revisited

- Let X be the number of pairwise collisions after k rolls
- What is $E[X]$?

Birthday Paradox, Revisited

- Let X be the number of pairwise collisions after k rolls

$$\begin{aligned} E[X] &= E[X_1 + \cdots + X_k] \\ &= E[X_1] + \cdots + E[X_k] \\ &= \frac{0}{n} + \cdots + \frac{k-1}{n} \\ &= \frac{k(k-1)}{2n} \end{aligned}$$

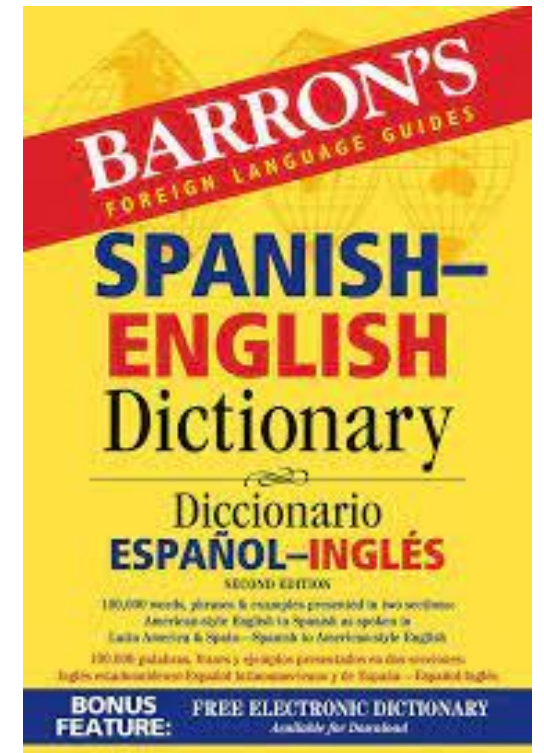
Birthday Paradox, Revisited

- $E[X] = \frac{k(k-1)}{2n}$
- $\frac{(k-1)^2}{2n} \leq E[X] \leq \frac{k^2}{2n}$
- $k = 2\sqrt{n} + 1$ implies $E[X] \geq 1$
- $k = \frac{\sqrt{n}}{2}$ implies $E[X] \leq \frac{1}{4}$

Case Study

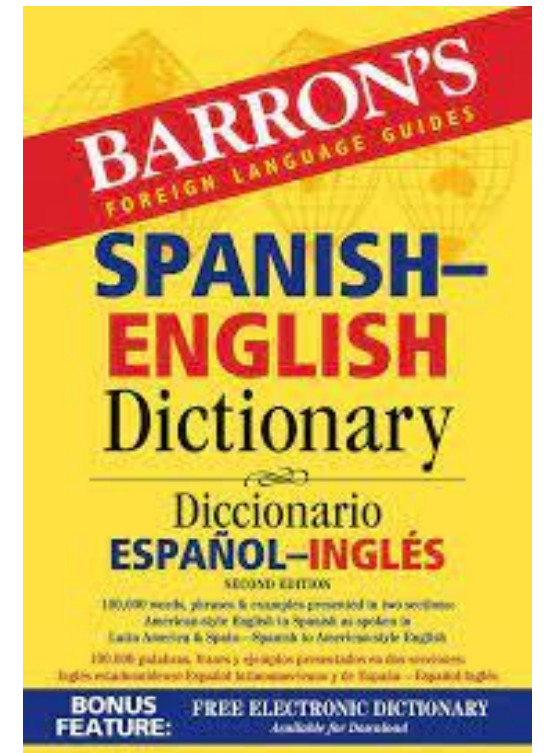
- We use the app for $k = 1000$ times and count the number of pairwise duplicates

- If the database contains *1 million words*, the expected number of pairwise duplicates is $E[X] = \frac{k(k-1)}{2n} < 0.5$



Case Study

- If the database contains *1 million words*, the expected number of pairwise duplicates is $E[X] = \frac{k(k-1)}{2n} < 0.5$
- ...We see **20** duplicates
- We think the claim is incorrect, but how can we be sure?



Concentration Inequalities

- Concentration inequalities bound the probability that a random variable is “far away” from its expectation
- Often used in understanding the performance of statistical tests, the behavior of data sampled from various distributions, and for our purposes, the guarantees of randomized algorithms

Markov's Inequality

- Let $X \geq 0$ be a non-negative random variable. Then for any $t > 0$:

$$\Pr[X \geq t \cdot E[X]] \leq \frac{1}{t}$$

Proof of Markov's Inequality

- Let $X \geq 0$ be a non-negative random variable. Then for any $t > 0$:

$$\begin{aligned} \mathbb{E}[X] &= \sum_{x \in \Omega} \Pr[X = x] \cdot x \\ &= \sum_{x \geq t \cdot \mathbb{E}[X]} \Pr[X = x] \cdot x + \sum_{x < t \cdot \mathbb{E}[X]} \Pr[X = x] \cdot x \\ &\geq \sum_{x \geq t \cdot \mathbb{E}[X]} \Pr[X = x] \cdot x \\ &\geq t \cdot \mathbb{E}[X] \sum_{x \geq t \cdot \mathbb{E}[X]} \Pr[X = x] \\ &= t \cdot \mathbb{E}[X] \cdot \Pr[X \geq t \cdot \mathbb{E}[X]] \end{aligned}$$

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Birthday Paradox, Revisited

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Birthday Paradox, Revisited

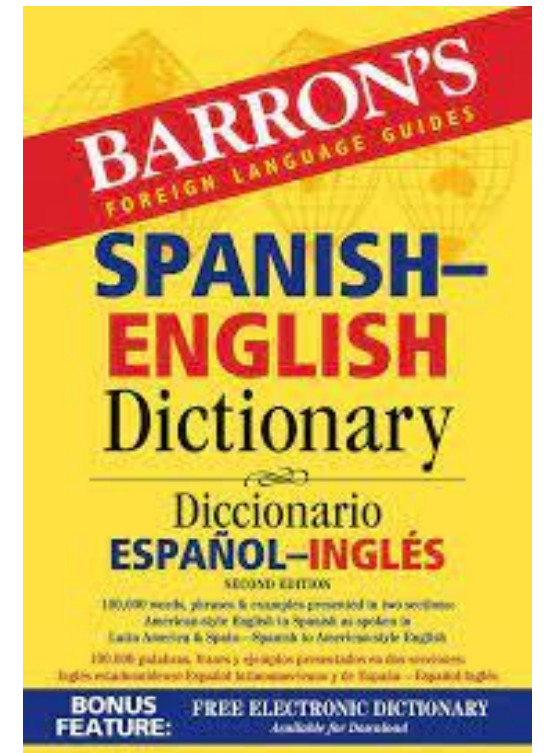
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Birthday Paradox, Revisited

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- $\frac{(k-1)^2}{2n} \leq E[X] \leq \frac{k^2}{2n}$
- $k = 2\sqrt{n} + 1$ implies $E[X] \geq 1$
- $k = \frac{\sqrt{n}}{2}$ implies $E[X] \leq \frac{1}{4}$, and by Markov's inequality, $\Pr[X \geq 1] \leq \frac{1}{4}$

Case Study

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Case Study

- If the database contains *1 million words*, the expected number of pairwise duplicates is $E[X] = \frac{k(k-1)}{2n} < 0.5$
- ...We see **20** duplicates
- $\Pr[X \geq 20] \leq \frac{1}{40}$

